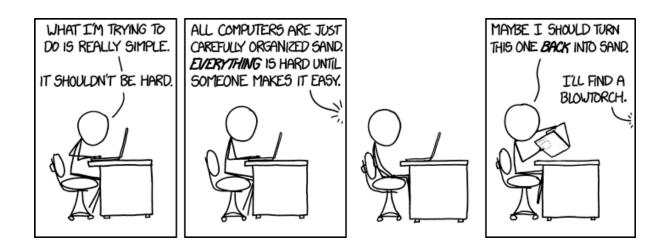
CITS3001 Mid-semester Test 2014

Fifty minutes

Answer all four questions

Total marks 60



When showing the operation of an algorithm, include enough detail to make it clear that you understand how *the algorithm* solves the problem.

Question 1: the Boyer-Moore algorithm for pattern-matching (15 marks)

Briefly describe the principles, operation, and performance issues of the Boyer-Moore algorithm for pattern-matching on strings.

Illustrate your answer using the DNA text and pattern in Fig. 1. Note that for DNA sequences, the alphabet is ACGT, i.e. only four characters.

T = TAGTCGTGCGCGAAATTCGTGCGCGA

P = TCGTGCGCG

Fig. 1: DNA text and pattern for Question 1.

Make sure that you show

- the complete pre-calculated tables for all heuristics;
- the shifts that are considered by the algorithm, the shifts that are skipped, and why; and
- all shifts that are returned as successful matches.

Question 2: the Travelling Salesman Problem

Briefly describe the principles, operation, and performance issues of any *one* commonly-used approximation algorithm for the Travelling Salesman Problem.

	Α	B	С	D	Ε
Α	—	40	10	20	12
В	40	—	35	50	32
С	10	35	—	30	16
D	20	50	30	—	19
E	12	32	16	19	-

Illustrate your answer using the set of cities in Fig. 2.

Make sure that you show all relevant operational details of your chosen algorithm.

Question 3: Iterative-deepening depth-first search

Briefly describe the principles, operation, and performance issues of iterative deepening depth-first search.

Illustrate your answer using the road map of Nerdvana in Fig. 3. Find the route from Geek to Nerd that goes through the fewest intermediate cities.

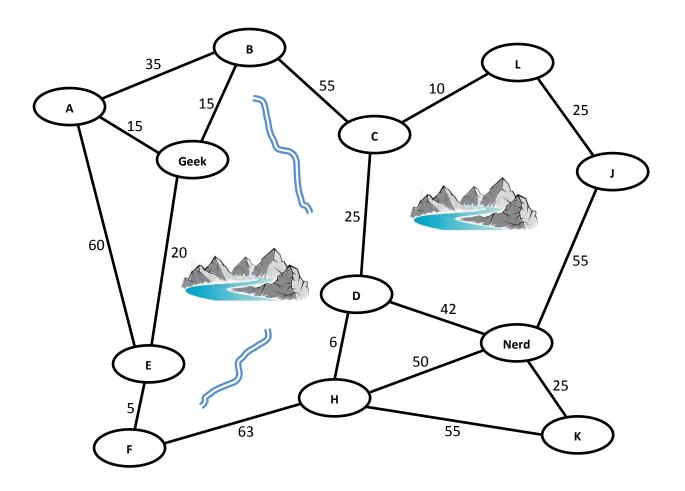


Fig. 3: Road map of Nerdvana for Questions 3 and 4.

Make sure that you show which nodes are expanded, in what order, and why.

Note that within one iteration, when you have expanded a node *X* at depth *k*, there is no need to expand *X* again at any depth $j \ge k$.

Question 4: *A**

Briefly describe the principles, operation, and performance issues of A^* .

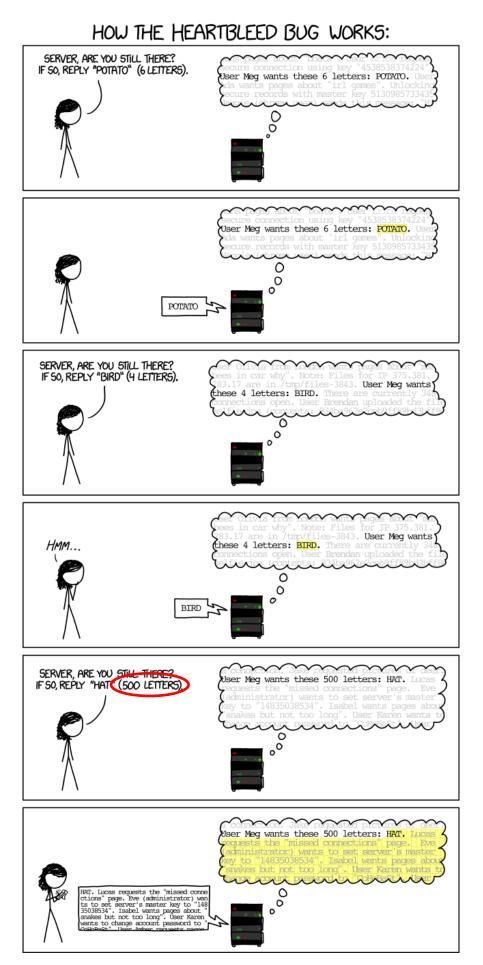
Illustrate your answer using the road map of Nerdvana in Fig. 3 and the table of straight-line distances in Fig. 4. Find the route from Geek to Nerd that requires the least total distance.

Geek	95	F	100
Α	125	Η	45
B	100	J	50
С	60	K	25
D	40	L	70
Ε	90	Nerd	0

Fig. 4: Straight-line distance from each city to Nerd for Question 4.

Make sure that you show which nodes are expanded, in what order, and why.

Note that when you have expanded a node X at g(X) = k, there is no need to expand X again at any $g(X) \ge k$.



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Solutions

Question 1

Need to mention:

- same basic system as the naïve algorithm;
- matches pattern right-to-left;
- uses the bad character heuristic;
- and the good-suffix heuristic;
- beats linear in good cases (or avoids some characters completely).

T = TAGTCGTGCGCGAAATTCGTGCGCGA

P = TCGTGCGCG

x[y] below for $\gamma(j)$ means x is the basic shift (P. 19), y is the advanced (P. 21).

b	λ(b)	j	γ(j)
Α	0	8	2[4]
С	8	7	2[6]
G	9	6	2
Т	4	≤5	9

Matching occurs from right-to-left at each shift.

T[9] = C, so the suggested shifts are 9-8 = 1 and 1

T[8–10] match, so the suggested shifts are 6-4 = 2 and 2

T[4–12] match, so the suggested shifts are 9 and 9

T[21] matches, so the suggested shifts are 8-4 = 4 and 2[4]

T[17–25] match, so the suggested shifts are 9 and 9 and we are done

Question 2

Various

Question 3

Need to mention five of:

- sequence of searches;
- each one is depth-limited;
- the limit increases in each iteration;
- space is linear in depth;
- optimal and complete;
- some nodes are inspected multiple times.

X* indicates that Node X has already been expanded.

```
Limit = 0, Geek is not a goal
```

Limit = 1, G(ABE)

```
Limit = 2, G(
```

A(BGE)

B(AGC)

E(AGF))

```
Limit = 3, G(
```

```
A(B(AGC)G*E(AGF))
```

B(A*G*C(BDL))

```
E(A*G*F(EH)))
```

Limit = 4, G(

 $A(B(A^*G^*C(BDL))G^*E(A^*G^*F(EH)))$

```
B(A*G*C(B*D(Nerd)L))
```

E)

So the solution with the fewest cities is GBCDN. (GEFHN is also acceptable.)

Question 4

Need to mention:

- maintains a pool/queue of unexpanded nodes, from the start node;
- lowest-cost node is selected next;
- cost of *n* is actual cost of start-to-*n* plus estimated cost of *n*-to-goal;
- provably optimal with admissible heuristic;
- good space behaviour requires *SMA**.

С	g(c)	h(c)	f(c)
Geek	0	95	95
EvenandC			•

Expand G

g(c)	h(c)	f(c)
20	90	110
15	100	115
15	125	140
	20 15	20 90 15 100

Expand E, ignore G and A, sort by f

g(c)	h(c)	f(c)
15	100	115
25	100	125
15	125	140
	15 25	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$

Expand B, ignore G and A, sort by f

С	g(c)	h(c)	f(c)
F	25	100	125
C	70	60	130
A	15	125	140
D 1 D		1 C	

Expand F, ignore E, sort by f

С	g(c)	h(c)	f(c)
C	70	60	130
Н	88	45	133
A	15	125	140

Expand C, ignore B, sort by f

С	g(c)	h(c)	f(c)
Н	88	45	133
D	95	40	135
A	15	125	140
L	80	70	150

С	g(c)	h(c)	f(c)
D	94	40	134
Ν	138	0	138
A	15	125	140
L	80	70	150
K	143	25	168

Expand H, ignore F, replace D, sort by f

Expand D, ignore C and H, replace N, sort by f

С	g(c)	h(c)	f(c)
Ν	136	0	136
А	15	125	140
L	80	70	150
K	143	25	168

So the solution with the least distance is GEFHDN.